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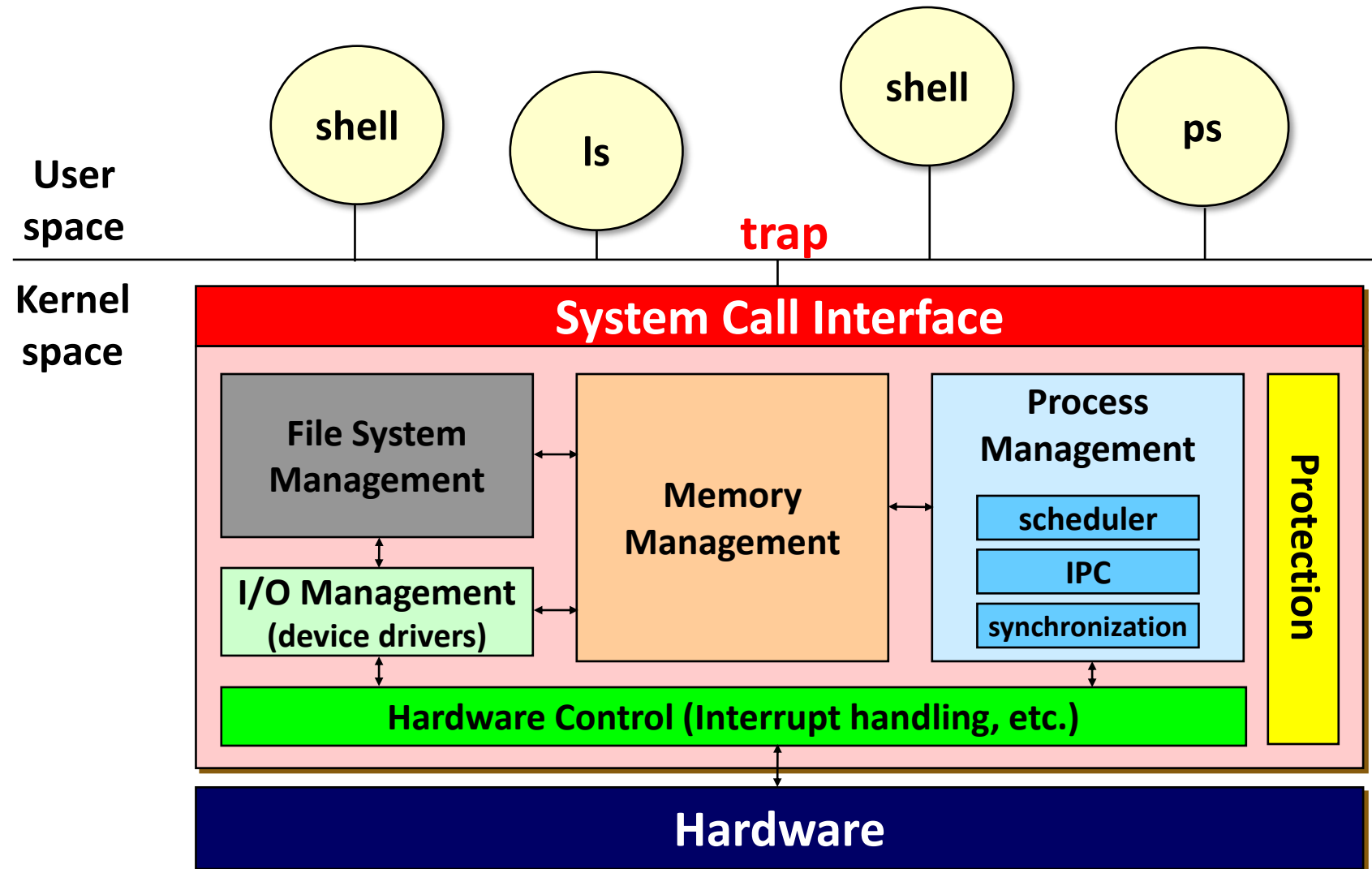
Systems Software &
Architecture Lab.
Seoul National University

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Processes



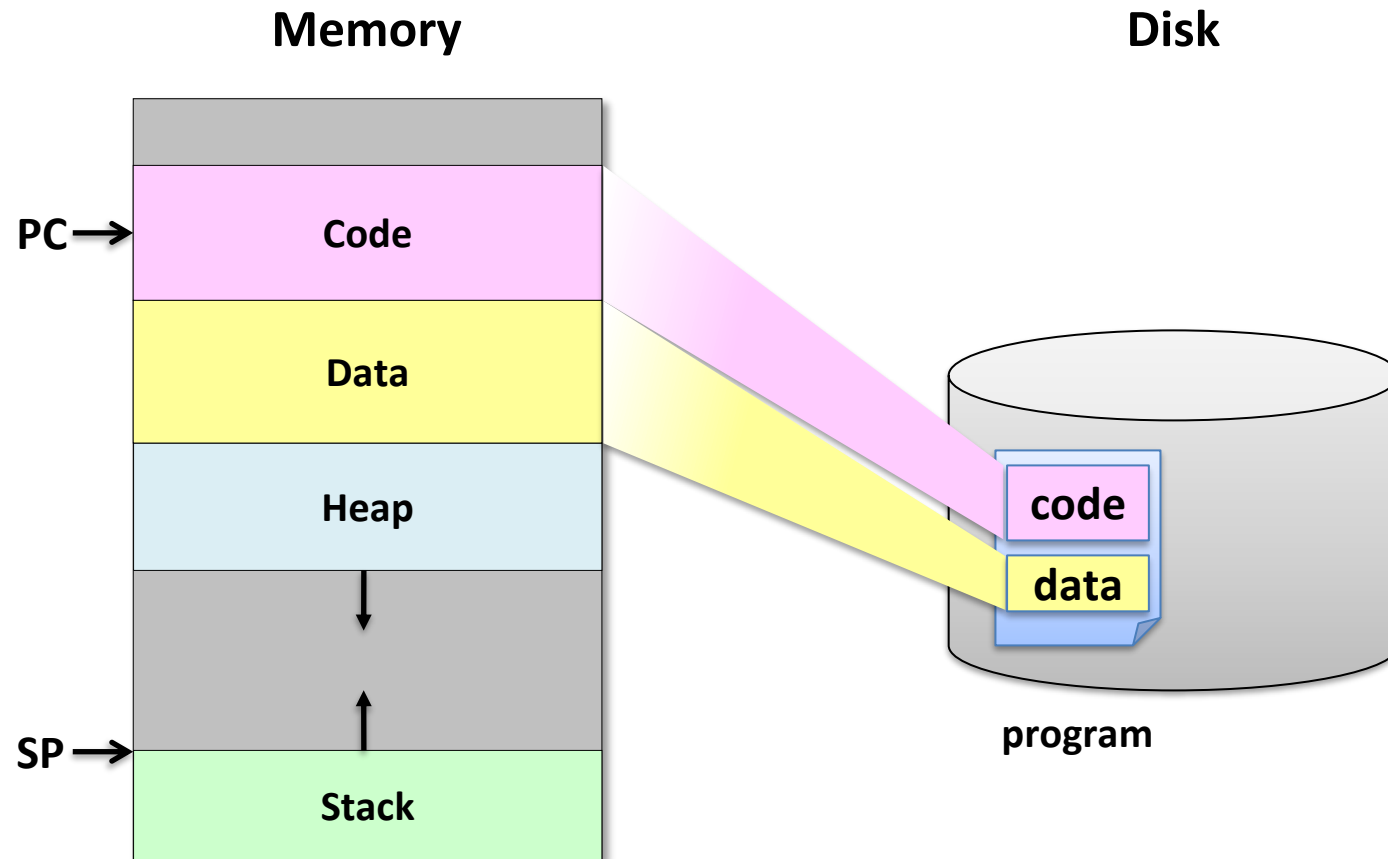
OS Internals



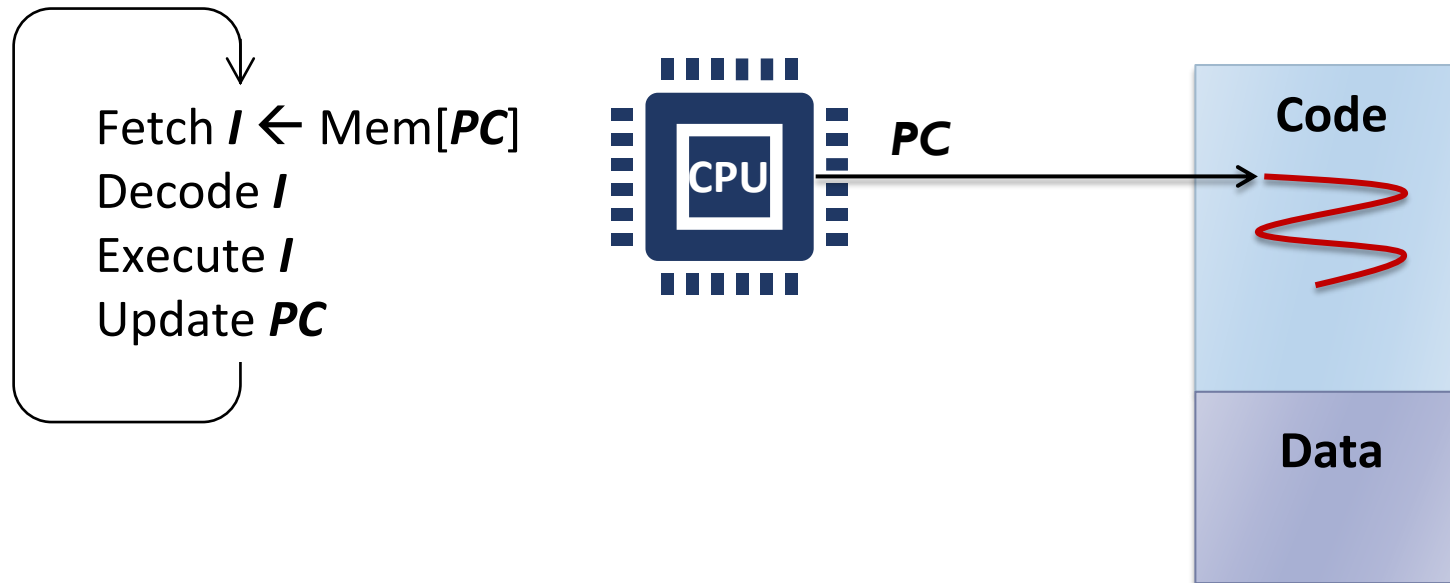
What is a Process?

- A(An) _____ of a program in execution
- Java analogy:
 - Class → “program” (static)
 - Object → “process” (dynamic)
- The basic unit of protection
- A process is identified using its process ID (PID)
- A process includes
 - CPU context (registers)
 - OS resources (address space, open files, etc.)
 - Other information (PID, state, owner, etc.)

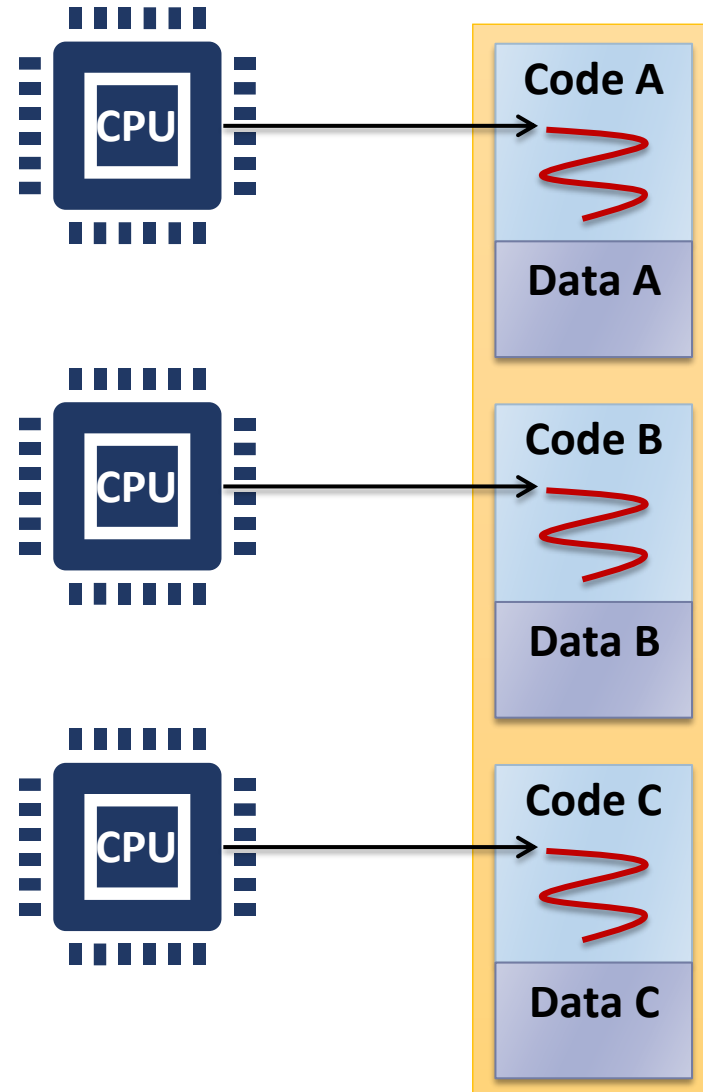
From Program to Process



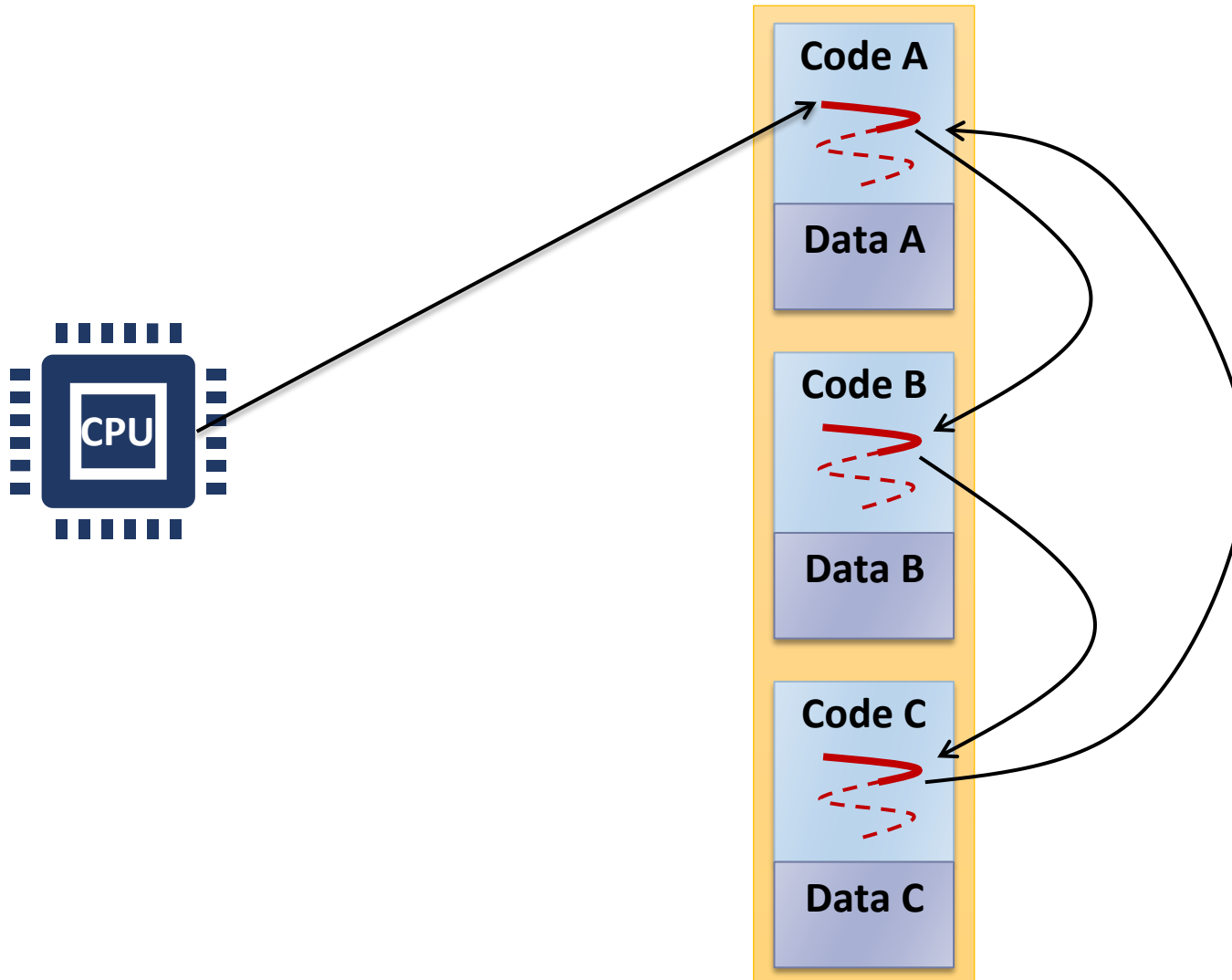
Running a Process



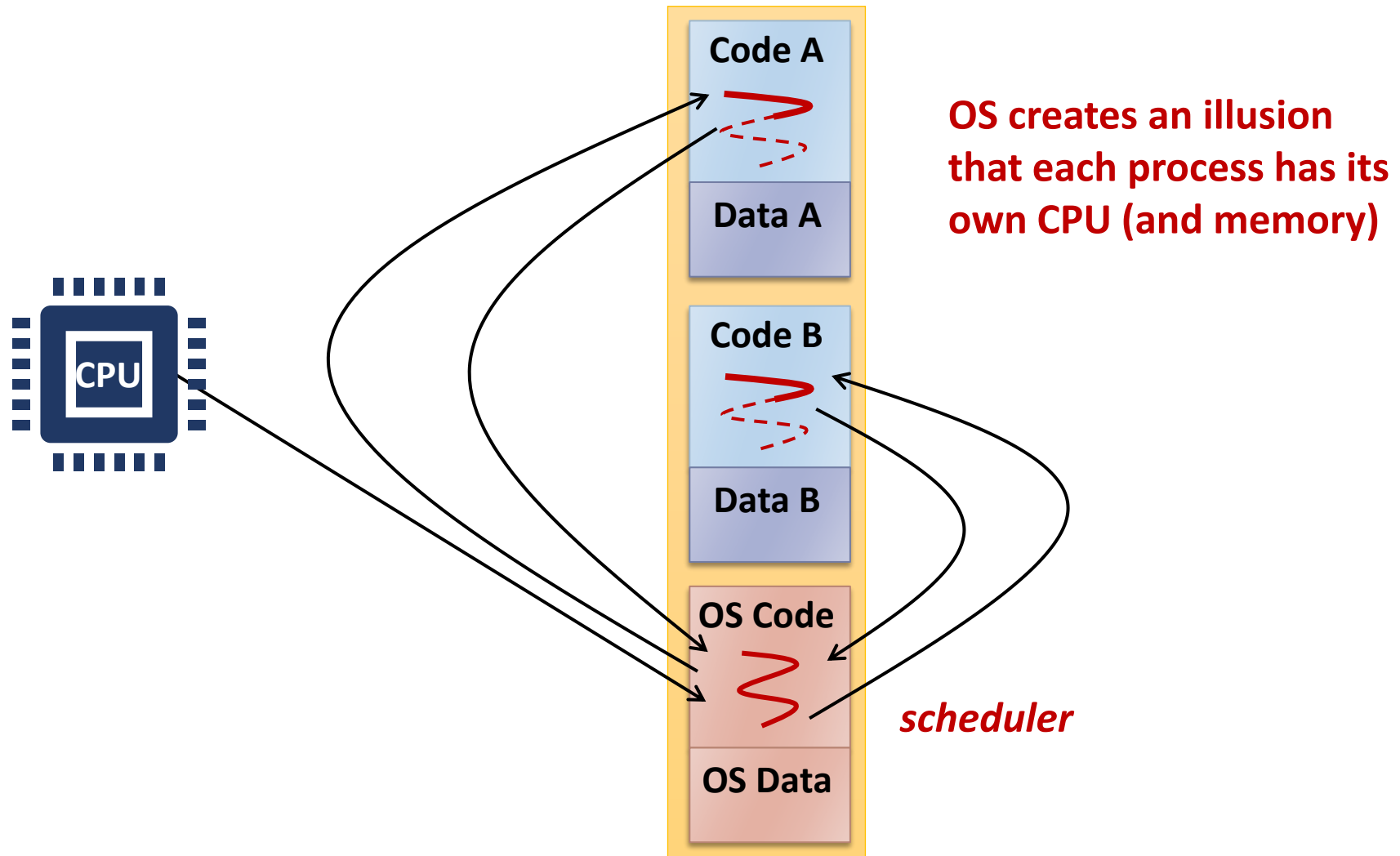
Running Multiple Processes



Interleaving Multiple Processes



Virtualizing the CPU



Example: Creating a Process

```
#include <sys/types.h>
#include <unistd.h>

int main() {
    int pid;

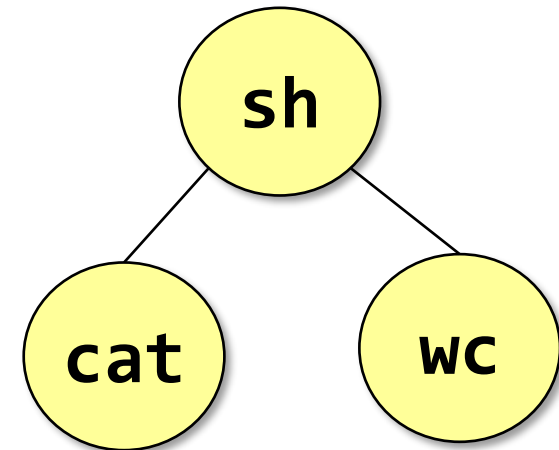
    if ((pid = fork()) == 0)
        printf ("Child of %d is %d\n", getppid(), getpid()); /* child */
    else
        printf ("I am %d. My child is %d\n", getpid(), pid); /* parent */
}
```

```
$ ./a.out
I am 31098. My child is 31099.
Child of 31098 is 31099.
$ ./a.out
Child of 31100 is 31101.
I am 31100. My child is 31101.
```

Process Hierarchy

- Parent-child relationship
 - One process can create another process
 - Unix calls the hierarchy a “process group”
 - Windows has no concept of process hierarchy
- Browsing a list of processes:
 - `ps` in Unix
 - Task Manager (`taskmgr`) in Windows

```
$ cat file1 | wc
```



Process Creation

▪ `fork()`

- Creates a new process cloning the parent process
 - Parent inherits most of resources and privileges: open files, UID, etc.
 - Child also duplicates the parent's address space
- Parent may either wait for the child to finish (using `wait()`), or it may continue in parallel
- Shells or GUIs use this system call internally
- Called once, returned twice

▪ `exec()`

- Replaces the current process image with a new program
- Windows: `CreateProcess()` = `fork()` + `exec()`
- Called once, never returns

Process Termination

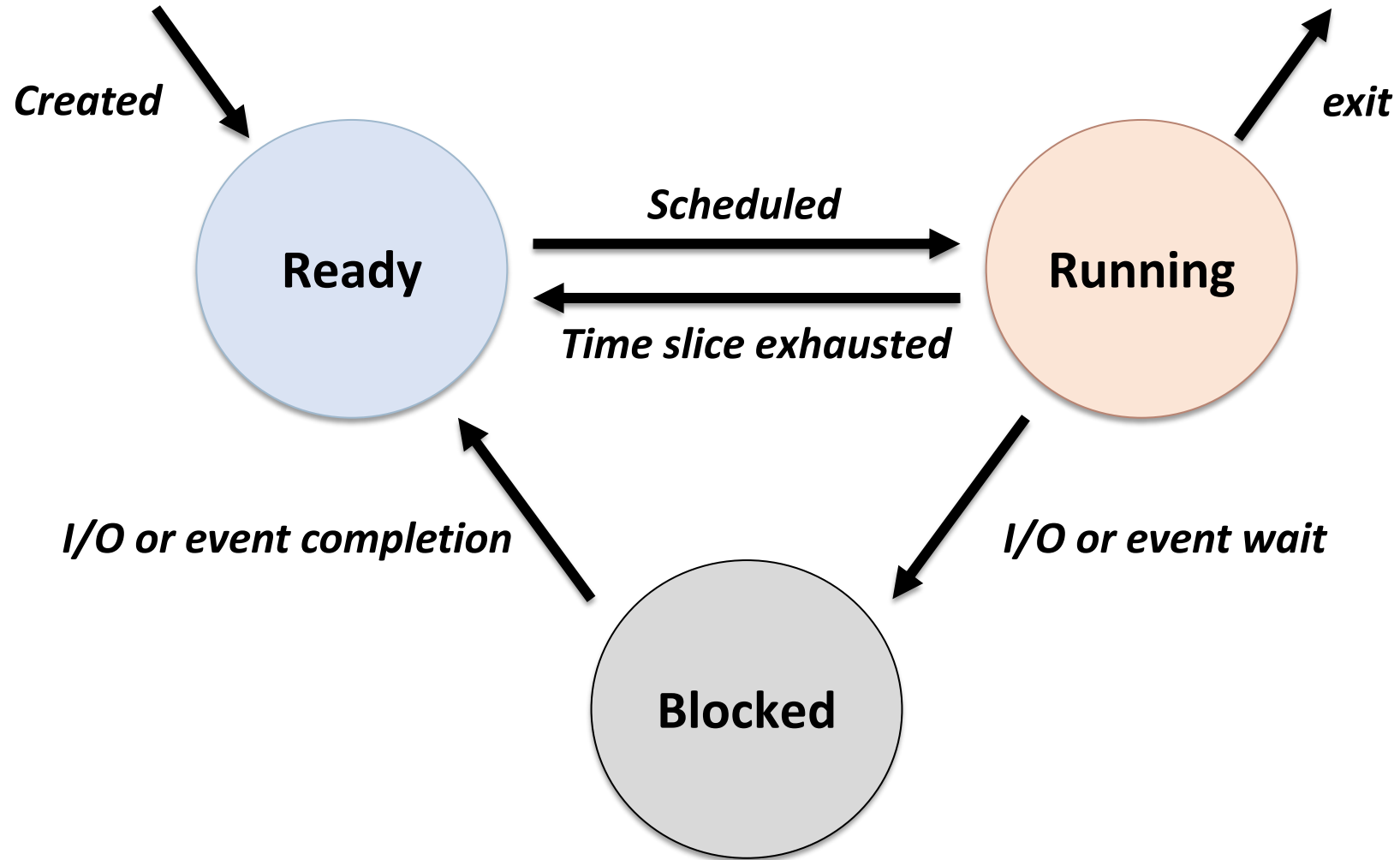
- Normal exit (voluntary)
- Error exit (voluntary)
- Fatal error (involuntary)
 - Segmentation fault – illegal memory access
 - Protection fault
 - Exceed allocated resources, etc.
- Killed by another process (involuntary)
 - By receiving a signal
- _____ process: terminated, but not removed

Simplified Shell

```
int main(void)
{
    char cmdline[MAXLINE];
    char *argv[MAXARGS];
    pid_t pid;
    int status;

    while (getcmd(cmdline, MAXLINE) >= 0) {
        parsecmd(cmdline, argv);
        if (!builtin_command(argv)) {
            if ((pid = fork()) == 0) {
                if (execv(argv[0], argv) < 0) {
                    printf("%s: command not found\n", argv[0]);
                    exit(0);
                }
            }
            waitpid(pid, &status, 0);
        }
    }
}
```

Process State Transitions



Processes

```

xterm
25041 ? S1 0:01 /usr/bin/epiphany
15124 ? Ss 0:01 /usr/sbin/nmbd -D
15126 ? Ss 0:00 /usr/sbin/smbd -D
15131 ? S 0:00 /usr/sbin/smbd -D
22930 ? S 0:10 /usr/sbin/smbd -D
3425 ? S 0:00 [pdflush]
20465 ? SMs 0:00 /usr/sbin/apache2 -k start
20479 ? SM 0:00 /usr/sbin/apache2 -k start
20480 ? SN 0:00 /usr/sbin/apache2 -k start
20481 ? SN 0:00 /usr/sbin/apache2 -k start
20482 ? SN 0:01 /usr/sbin/apache2 -k start
20483 ? SN 0:01 /usr/sbin/apache2 -k start
4762 ? SN 0:01 /usr/sbin/apache2 -k start
4952 ? SN 0:00 /usr/sbin/apache2 -k start
4953 ? SN 0:00 /usr/sbin/apache2 -k start
31647 ? SN 0:01 /usr/sbin/apache2 -k start
32071 ? SN 0:00 /usr/sbin/apache2 -k start
3708 ? Ss 0:00 sshd: jinsoo [priv]
3710 ? S 0:00 sshd: jinsoo@notty
3711 ? Ss 0:00 tcsh -c xterm
3716 ? S 0:00 xterm -g 80x30 -fg white -bg #003333 -sb -sl 5000 -cr
3717 pts/0 Ss+ 0:00 -csh
3934 ? Ss 0:00 sshd: jinsoo [priv]
3936 ? S 0:00 sshd: jinsoo@notty
3937 ? Ss 0:00 tcsh -c xterm
3942 ? S 0:00 xterm -g 80x30 -fg white -bg #003333 -sb -sl 5000 -cr
3943 pts/1 Ss 0:00 -csh
3981 ? Ss 0:00 imapd
3997 pts/1 R+ 0:00 ps ax
[oz:/user/jinsoo-3]

```

작업 관리자

파일(F) 옵션(O) 보기(V)

프로세스 성능 앱 기록 시작프로그램 사용자 세부 정보 서비스

이름	상태	5% CPU	19% 메모리	1% 디스크	0% 네트워크	0% GPU
> 작업 관리자		0.6%	36.2MB	0MB/s	0Mbps	0%
> Initech Client Manager Service...		0.4%	2.4MB	0MB/s	0Mbps	0%
> TUCTSystem.exe(32비트)		0.4%	2.8MB	0MB/s	0Mbps	0%
System		0.3%	0.1MB	0.1MB/s	0Mbps	0%
> 서비스 호스트: Windows Update		0.3%	15.0MB	0.1MB/s	0Mbps	0%
데스크톱 창 관리자		0.3%	74.1MB	0MB/s	0Mbps	0.1%
서비스 및 컨트롤러 응용 프로...		0.3%	6.1MB	0MB/s	0Mbps	0%
> 서비스 호스트: Windows Mana...		0.3%	9.0MB	0MB/s	0Mbps	0%
Antimalware Service Executable		0.2%	181.7MB	0.1MB/s	0Mbps	0%
Slack		0.2%	45.6MB	0MB/s	0Mbps	0%
> MagicLine4NXServices(32비트)		0.1%	7.4MB	0MB/s	0Mbps	0%
> CrossEX Live Checker(32비트)		0.1%	1.4MB	0MB/s	0Mbps	0%
> Microsoft Edge(23)		0.1%	1,311.4MB	0MB/s	0Mbps	0%
> Google Chrome(4)		0.1%	109.8MB	0MB/s	0Mbps	0%
> Windows 탐색기(2)		0.1%	84.4MB	0MB/s	0Mbps	0%
ASDF Service Application		0.1%	4.9MB	0MB/s	0Mbps	0%
WMI Provider Host		0.1%	30.4MB	0MB/s	0Mbps	0%
Dell Display Manager(32비트)		0.1%	2.4MB	0MB/s	0Mbps	0%
> Spooler SubSystem App		0.1%	5.8MB	0MB/s	0Mbps	0%
> 서비스 호스트: Network List Ser...		0.1%	2.8MB	0MB/s	0Mbps	0%
> Interezen Service Program(32...		0.1%	1.5MB	0MB/s	0Mbps	0%
WMI Provider Host		0.1%	2.7MB	0MB/s	0Mbps	0%

간단히(D) 작업 끝내기(E)

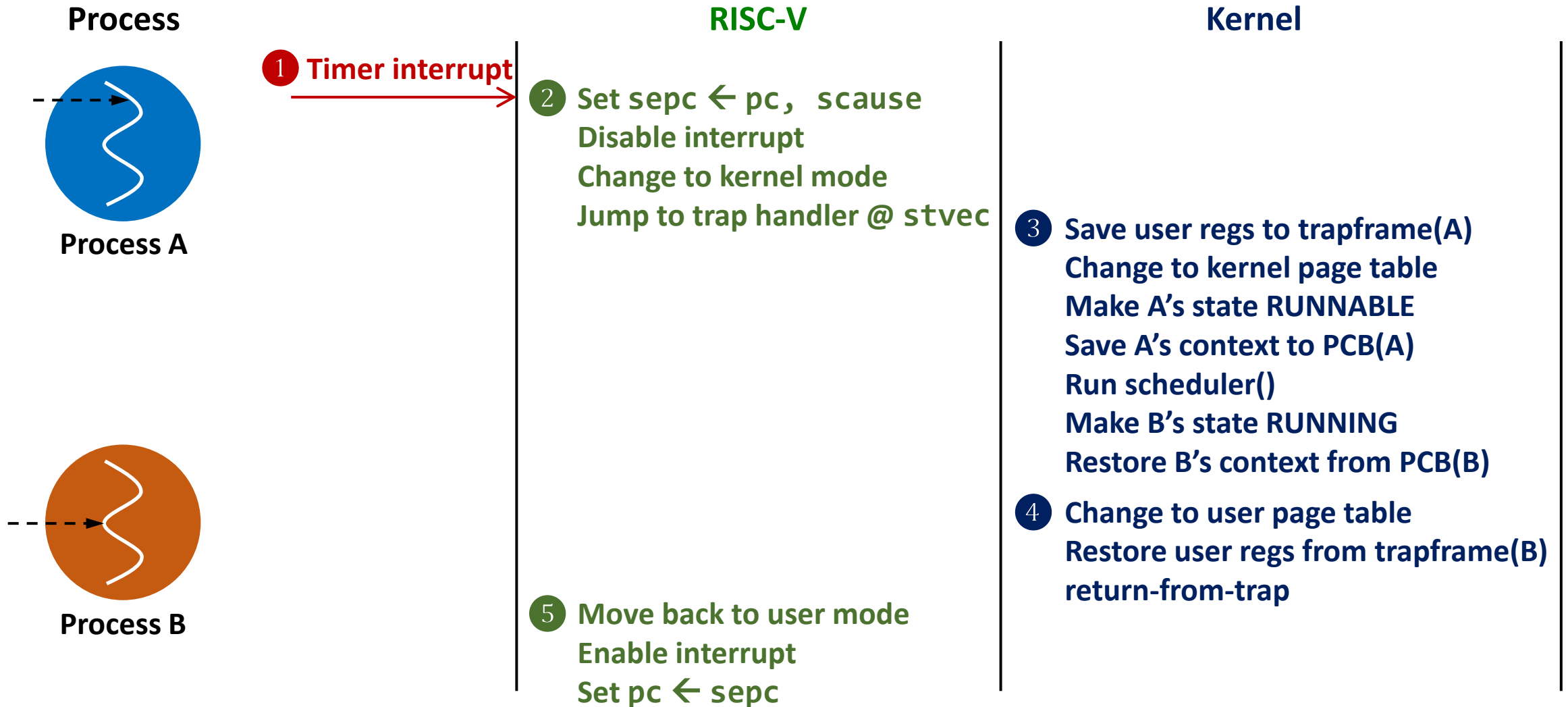
Implementing Processes

- PCB (Process Control Block) or Process Descriptor
 - Each PCB represents a process
 - Contains all the information about a process
 - CPU registers
 - PID, PPID, process group, priority, process state, signals
 - CPU scheduling information
 - Memory management information
 - Accounting information
 - File management information
 - I/O status information
 - Credentials
 - `struct task_struct` in Linux: 6592 bytes as of Linux 6.2.0
 - `struct proc` in xv6: 360 bytes

Context Switch

- The act of switching CPU from one process to another
- Administrative overhead
 - Saving and restoring registers and memory maps
 - Flushing and reloading the memory cache
 - Updating various tables and lists, etc.
- The overhead depends on hardware support
 - Multiple register sets in UltraSPARC
 - Advanced memory management techniques may require extra data to be switched with each context (e.g., page tables, TLB, etc.)
- 100s or 1000s of switches/sec typically

Performing Context Switch in xv6



Process State Queues

- The OS maintains a collection of queues that represent the state of all processes in the system
 - Ready queue (or run queue)
 - Wait queue(s): one queue for each type of event (device, timer, message, ...)
- Each PCB is queued onto a state queue according to its current state
 - As a process changes state, its PCB is migrated between the various queues

Implementing fork()

```
int fork()
```

- Creates and initializes a new PCB
- Creates and initializes a new address space
- Initializes the address space with a copy of the entire contents of the address space of the parent
- Initializes the kernel resources to point to the resources used by the parent (e.g., open files)
- Places the PCB on the ready queue
- Returns the child's PID to the parent, and zero to the child

Implementing `exec()`

```
int execv(char *prog, char *argv[])
```

- Stops the current process
- Loads the program “prog” into the process’s address space
- Initializes hardware context and “args” for the new program
- Places the PCB on the ready queue

- `exec()` does not create a new process
- What does it mean for `exec()` to return?

Policy vs. Mechanism

■ Policy

- *What* should be done?
- Policy decisions must be made for all resource allocation and scheduling problems
- e.g., What is the next process to run?

■ Mechanism

- *How* to do something?
- The tool for implementing a set of policies
- e.g., How to make multiple processes run at once?

Separating Policy from Mechanism

- A key principle in operating system design
- Policies are likely to change depending on workloads and also across places or over time
- A general mechanism, separated from policy, is more desirable
- Allows to build a more modular OS
- Enables extensible systems – User-specific policies?