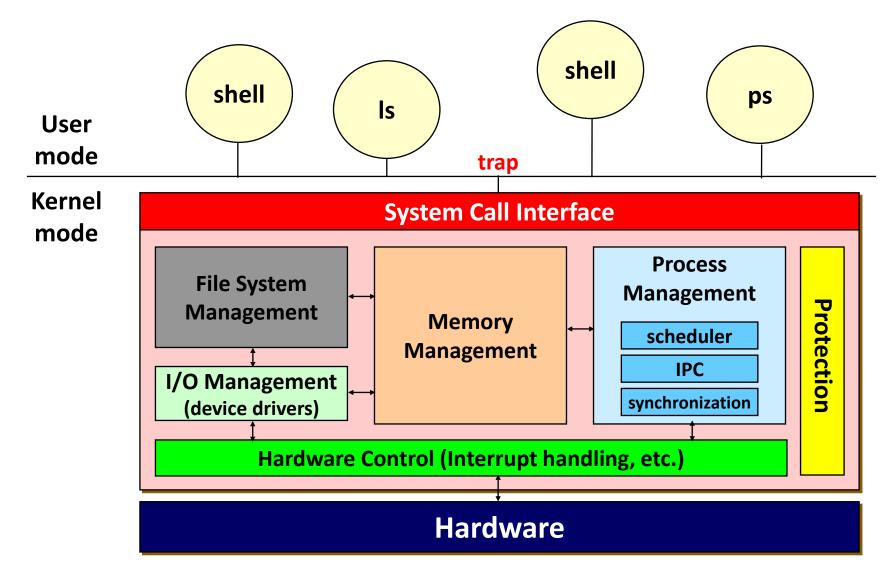
Jin-Soo Kim (jinsoo.kim@snu.ac.kr) Systems Software & Architecture Lab. Seoul National University

Spring 2024

Introduction to Operating Systems



Operating System Internals



OS: Application View

- OS provides an execution environment for running programs

OS provides a(an) ______ view of the underlying computer system

- What are the correct abstractions?
- How much of hardware should be exposed?
- Typical OS abstractions
 - Processors \rightarrow Processes, Threads
 - Memory \rightarrow Address space (virtual memory)
 - Storage \rightarrow Volumes, Directories, Files
 - I/O Devices \rightarrow Files (+ ioctls)
 - Networks \rightarrow Files (sockets, pipes, ...)



OS: System View

- OS manages various resources of a computer system
- Sharing

- Fairness
- Efficiency

CPU
Memory
I/O devices
Queues
Energy
...

Resources

OS: Implementation View

OS is highly-concurrent, software Two kinds of events trap • System calls • Interrupts System call Interrupts Hardware

Unix Features

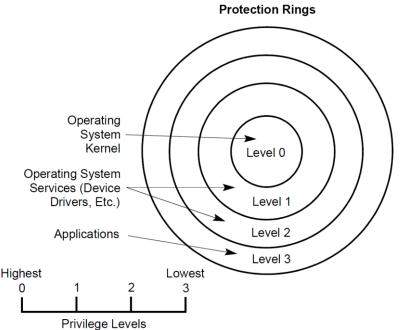
- Process control
 - fork(), exec(), wait(), exit()
 - Pipes for inter-process communication (IPC)
- Hierarchical file systems
 - Special files: uniform I/O, naming, and protection
 - Removable file systems via mount/umount
 - i-node
- Signals
- Shells
 - Standard I/O and I/O redirection, filters
 - Shell scripts

Architectural Support for OS (1)

- CPU modes of operation: kernel vs. user
 - 4 levels in x86: Ring 0 > 1 > 2 > 3
 - 3 levels in RISC-V: Machine > Supervisor > User
- Protected or ______ instructions
 - Direct I/O access (e.g., in/out instructions in x86)
 - Accessing system registers
 - Memory state management

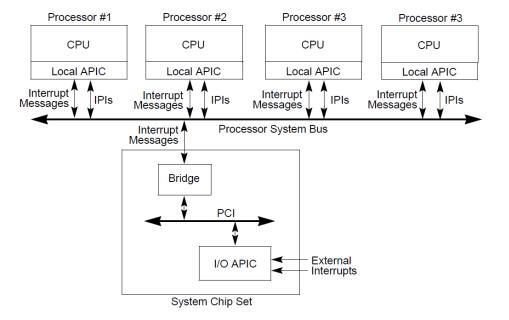
•

. . .



Architectural Support for OS (2)

- Interrupts
 - Generated by hardware devices
 - External interrupts vs. IPIs
 - Asynchronous
- Exceptions
 - Generated by software executing instructions
 - Faults (unintentional, but possibly recoverable): page faults, protection faults, ...
 - Traps (intentional): syscall instruction in x86_64 or ecall instruction in RISC-V
 - Aborts (unintentional and unrecoverable): parity error, machine error, ...
 - Synchronous
 - Exception handling is logically same as interrupt handling



Architectural Support for OS (3)

- Memory protection
 - Segmentation
 - Paging
- Timer

. . .

- DMA (Direct Memory Access)
- Atomic instructions
 - Atomic inc/dec
 - Test-and-Set
 - Compare-and-Swap
 - LL (Load Locked) & SC (Store Conditional)

System Calls

System Calls

OS defines a set of system calls

- Programming interface to the services provided by OS
- OS protects the system by rejecting illegal requests
- OS may impose a quota on a certain resource
- OS may consider fairness while sharing a resource
- A system call is a _____ procedure call
 - System call routines are in the OS code
 - Executed in the kernel mode
 - On entry, user mode \rightarrow kernel mode switch
 - On exit, CPU mode is changed back to the user mode

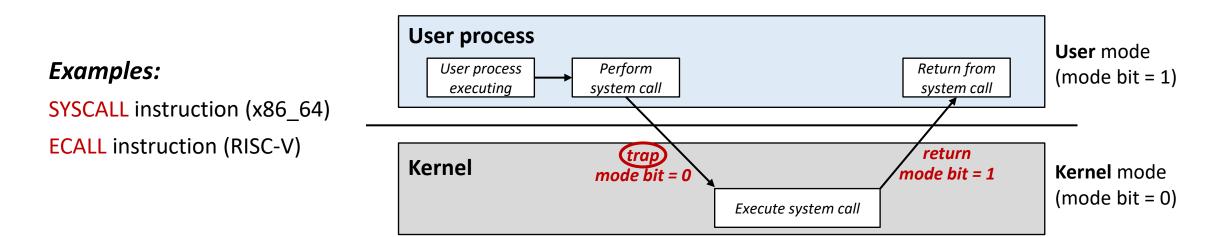
System Calls Example

POSIX vs.Win32

Category	POSIX	Win32	Description		
	fork	CreateProcess	Create a new process (CreateProcess = fork + exec)		
Process	waitpid	WaitForSingleObject	Wait for a process to exit		
	execve	(none)	Execute a new program		
Management	exit	ExitProcess	Terminate execution		
	kill	(none)	Send a signal		
	open	CreateFile	Create a file or open an existing file		
	close	CloseHandle	Close a file		
File	read	ReadFile	Read data from a file		
	write	WriteFile	Write data to a file		
Management	lseek	SetFilePointer	Move the file pointer		
	stat	GetFileAttibutesEx	Get various file attributes		
	chmod	(none)	Change the file access permission		
	mkdir	CreateDirectory	Create a new directory		
	rmdir	RemoveDirectory	Remove an empty directory		
File System	link	(none)	Make a link to a file		
Management	unlink	DeleteFile	Destroy an existing file		
	chdir	SetCurrentDirectory	Change the current working directory		
	mount	(none)	Mount a file system		

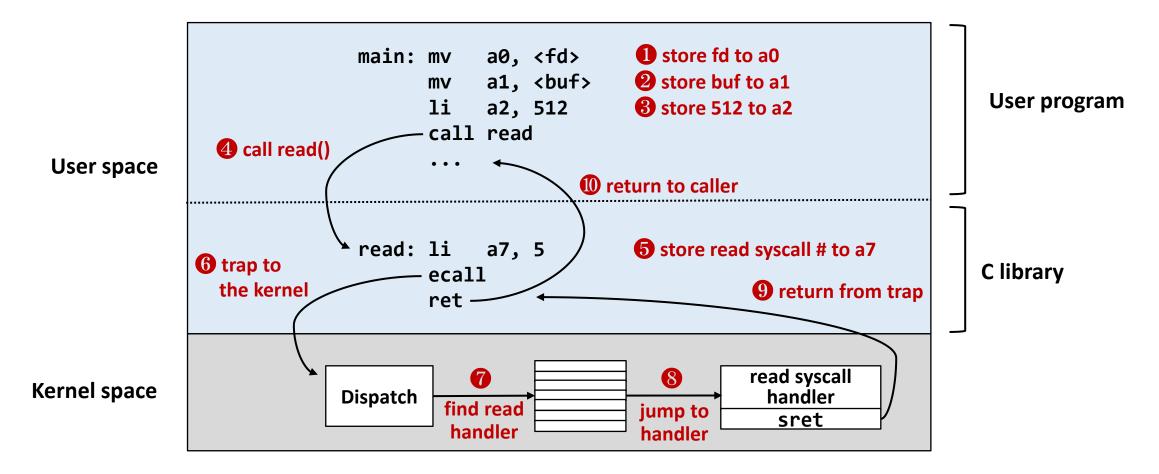
OS Trap

- There must be a special "trap" instruction that:
 - Causes an exception, which invokes a kernel handler
 - Passes a parameter indicating which system call to invoke
 - Saves caller's state (registers, mode bits)
 - Returns to user mode when done with restoring its state
 - OS must verify caller's parameters (e.g., pointers)

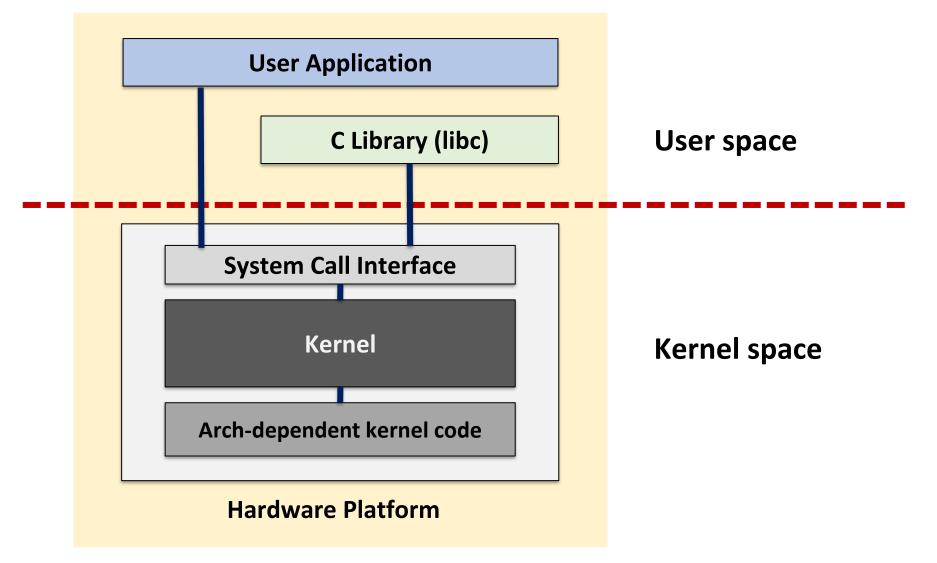


Implementing System Calls

count = read(fd, buf, 512);



Typical (Monolithic) OS Structure



Trap Instructions in x86

- int 0x80(+iret)
 - "Software interrupt"
 - A legacy way to invoke a system call (used for 32-bit mode)
 - Slow: use the same mechanism as interrupts
- sysenter (+ sysexit)
 - A new, fast instruction to invoke a system call in 32-bit mode
 - Introduced by Intel (not available in 64-bit mode on AMD CPUs)
- syscall (+ sysret)
 - Similar to **sysenter**, but used in **64-bit mode**
 - Introduced by AMD (not available in 32-bit mode on Intel CPUs)



SYSCALL—Fast System Call

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 05	SYSCALL	ZO	Valid	Invalid	Fast call to privilege level 0 system procedures.

- $RCX \leftarrow RIP$
- RIP ← MSR_LSTAR
- R11 \leftarrow RFLAGS
- •
- Initialize CS and SS from MSR_STAR
- Set CPL(Current Privilege Level) to 0

Using a System Call

Example:getpid()

```
#include <sys/types.h>
#include <unistd.h>
int main(void)
{
   return getpid();
}
```

Invoking a System Call

System call number for getpid(): 0x27 (= 39)

00000000004005a0 <main>: 4005a0: e9 fb 84 04 00 4005a5: 66 2e 0f 1f 84 00 00 4005ac: 00 00 00</main>	jmpq 448aa0 <getpid> nopw %cs:0x0(%rax,%rax,1)</getpid>	
4005af: 90	nop	
000000000448aa0 <getpid>:</getpid>		

 448aa0
 <__getpla>:

 448aa0:
 b8 27 00 00 00
 mov
 \$0x27,%eax

 448aa5:
 0f 05
 syscall

 448aa7:
 c3
 retq

 448aa8:
 0f 1f 84 00 00 00
 nopl
 0x0(%rax,%rax,1)

 448aaf:
 00
 00
 00
 00

In C library

Numbering System Calls

@ arch/x86/entry/syscalls/syscall_64.tbl

#	64-bit system call numbers and entry vectors	
#	The format is:	
	<number> <abi> <name> <entry point=""></entry></name></abi></number>	
#	Thex64_sys_*() stubs are created on-the-fly for sys_*() system calls	
#		
#	The abi is "common", "64" or "x32" for this file.	
# 0	common readx64_sys_read	
1	common writex64_sys_write	
2	common openx64_sys_open	
3	common closex64_sys_close	
4	common statx64_sys_newstat	
3	7 common alarmx64_sys_alarm	
3	8 common setitimerx64_sys_setitimer	
3		
4	0 common sendfilex64_sys_sendfile64	

Setting up the System Call Entry

@ arch/x86/kernel/cpu/common.c

```
void syscall_init(void)
{
```

```
wrmsr(MSR_STAR, 0, (__USER32_CS << 16) | __KERNEL_CS);
wrmsrl(MSR_LSTAR, (unsigned long)entry_SYSCALL_64);
```

Entering the Kernel

@ arch/x86/entry/entry_64.S

SYM_CODE_START(entry_SYSCALL_64) UNWIND_HINT_EMPTY

swapgs

```
/* tss.sp2 is scratch space. */
movq %rsp, PER_CPU_VAR(cpu_tss_rw + TSS_sp2)
SWITCH_TO_KERNEL_CR3 scratch_reg=%rsp
movq PER_CPU_VAR(cpu_current_top_of_stack), %rsp
```

```
/* Construct struct pt_regs on stack */
pushq $__USER_DS /* pt_regs->ss */
pushq PER_CPU_VAR(cpu_tss_rw + TSS_sp2) /* pt_regs->sp */
pushq %r11 /* pt_regs->flags */
pushq $__USER_CS /* pt_regs->cs */
pushq %rcx /* pt_regs->ip */
SYM_INNER_LABEL(entry_SYSCALL_64_after_hwframe, SYM_L_GLOBAL)
pushq %rax /* pt_regs->orig_ax */
```

PUSH_AND_CLEAR_REGS rax=\$-ENOSYS

```
/* IRQs are off. */
movq %rax, %rdi
movq %rsp, %rsi
call do_syscall_64 /* returns with IRQs disabled */
```

Jumping to the System Call Handler

@ arch/x86/entry/common.c

```
__visible void do_syscall_64(unsigned long nr, struct pt_regs *regs)
struct thread_info *ti;
 enter_from_user_mode();
 local_irq_enable();
 ti = current_thread_info();
 if (READ_ONCE(ti->flags) & _TIF_WORK_SYSCALL_ENTRY)
   nr = syscall_trace_enter(regs);
 if (likely(nr < NR_syscalls)) {</pre>
   nr = array_index_nospec(nr, NR_syscalls);
   regs->ax = sys_call_table[nr](regs);
```

System Call Table

@ arch/x86/entry/syscall_64.c

```
#define __SYSCALL_64(nr, sym) extern long __x64_##sym(const struct pt_regs *);
#include <asm/syscalls_64.h>
#undef __SYSCALL_64
#define __SYSCALL_64(nr, sym) [nr] = __x64_##sym,
asmlinkage const sys_call_ptr_t sys_call_table[__NR_syscall_max+1] = {
  /*
   * Smells like a compiler bug -- it doesn't work
   * when the & below is removed.
   */
  [0 ... __NR_syscall_max] = &__x64_sys_ni_syscall,
#include <asm/syscalls_64.h>
};
```

System Call Handlers

@ arch/x86/include/generated/asm/syscalls_64.h

__SYSCALL_COMMON(0, sys_read) __SYSCALL_COMMON(1, sys_write) __SYSCALL_COMMON(2, sy<u>s_open)</u> __SYSCALL_COMMON(3, sys_close) __SYSCALL_COMMON(4, sys_newstat) ___SYSCALL_COMMON(5, sys_newfstat) __SYSCALL_COMMON(6, sys_newlstat) __SYSCALL_COMMON(7, sys_poll) __SYSCALL_COMMON(8, sys_lseek) __SYSCALL_COMMON(9, sys_mmap) __SYSCALL_COMMON(10, sys_mprotect) __SYSCALL_COMMON(11, sys_munmap) __SYSCALL_COMMON(12, sys_brk) __SYSCALL_64(13, sys_rt_sigaction) __SYSCALL_COMMON(14, sys_rt_sigprocmask) __SYSCALL_64(15, sys_rt_sigreturn) __SYSCALL_64(16, sys_ioctl)

vDSO

- Virtual Dynamic Shared Object (@ arch/x86/entry/vdso/)
 - A small shared library exported by the kernel that is mapped into the address space of all user-space applications
 - Mapped to a different location every time (for security)
 - Used to accelerate the execution of certain read-only system calls ("virtual system calls") without entering the kernel
 - clock_gettime()
 - gettimeofday()
 - getcpu()
 - time()
 - clock_getres()
 - @ arch/x86/entry/vdso
 - \$ man vdso

vDSO Layout

<pre>\$ cat /proc/self/maps</pre>				
5611a0174000-5611a0176000	rp	00000000	08:05	265118
5611a0176000-5611a017b000	r-xp	00002000	08:05	265118
5611a017b000-5611a017e000	rp	00007000	08:05	265118
5611a017e000-5611a017f000	rp	00009000	08:05	265118
5611a017f000-5611a0180000	rw-p	0000a000	08:05	265118
5611a1ebb000-5611a1edc000	rw-p	00000000	00:00	0
7f9d19bdc000-7f9d19bfe000	rw-p	00000000	00:00	0
7f9d19bfe000-7f9d1ab21000	rp	00000000	08:05	269248
7f9d1ab21000-7f9d1ab46000	rp	00000000	08:05	274111
7f9d1ab46000-7f9d1acbe000	r-xp	00025000	08:05	274111
7f9d1acbe000-7f9d1ad08000	rp	0019d000	08:05	274111
7f9d1ad08000-7f9d1ad09000	р	001e7000	08:05	274111
7f9d1ad09000-7f9d1ad0c000	rp	001e7000	08:05	274111
7f9d1ad0c000-7f9d1ad0f000	rw-p	001ea000	08:05	274111
7f9d1ad0f000-7f9d1ad15000	rw-p	00000000	00:00	0
7f9d1ad24000-7f9d1ad25000	rp	00000000	08:05	273898
7f9d1ad25000-7f9d1ad48000	r-xp	00001000	08:05	273898
7f9d1ad48000-7f9d1ad50000	rp	00024000	08:05	273898
7f9d1ad51000-7f9d1ad52000	rp	0002c000	08:05	273898
7f9d1ad52000-7f9d1ad53000	rw-p	0002d000	08:05	273898
7f9d1ad53000-7f9d1ad54000	rw-p	00000000	00:00	0
7ffe8b35b000-7ffe8b37c000	rw-p	00000000	00:00	0
7ffe8b3f3000-7ffe8b3f6000	rp	00000000	00:00	0
7ffe8b3f6000-7ffe8b3f7000	r-xp	00000000	00:00	0
fffffffff60000-fffffff	ff6010	000xp	000000	00:00

/usr/bin/cat /usr/bin/cat /usr/bin/cat /usr/bin/cat [heap]

/usr/lib/locale/locale-archive /usr/lib/x86_64-linux-gnu/libc-2.31.so /usr/lib/x86_64-linux-gnu/libc-2.31.so /usr/lib/x86_64-linux-gnu/libc-2.31.so /usr/lib/x86_64-linux-gnu/libc-2.31.so /usr/lib/x86_64-linux-gnu/libc-2.31.so

/usr/lib/x86_64-linux-gnu/ld-2.31.so
/usr/lib/x86_64-linux-gnu/ld-2.31.so
/usr/lib/x86_64-linux-gnu/ld-2.31.so
/usr/lib/x86_64-linux-gnu/ld-2.31.so

[stack] [vvar] [vdso] [vsyscall]

0

vvar and vsyscall

- vvar
 - Mapped just before the vdso page
 - Contains data accessed by virtual system calls
 - Kernel periodically updates the values (if necessary)
 - User-space application can only read the values

vsyscall

- A legacy ABI for virtual system calls
- Mapped to the fixed user-space address
- Not recommended to use



- A generic library function that performs the specified system call
 - Symbolic constants for system call numbers are specified in <sys/syscall.h>
 - Useful when you add a new system call that has no wrapper function in the C library

```
#include <stdio.h>
#include <unistd.h>
#include <sys/syscall.h>
int main(void)
{
    printf("%ld\n", syscall(__NR_getpid));
    printf("%ld\n", syscall(SYS_getpid));
}
```