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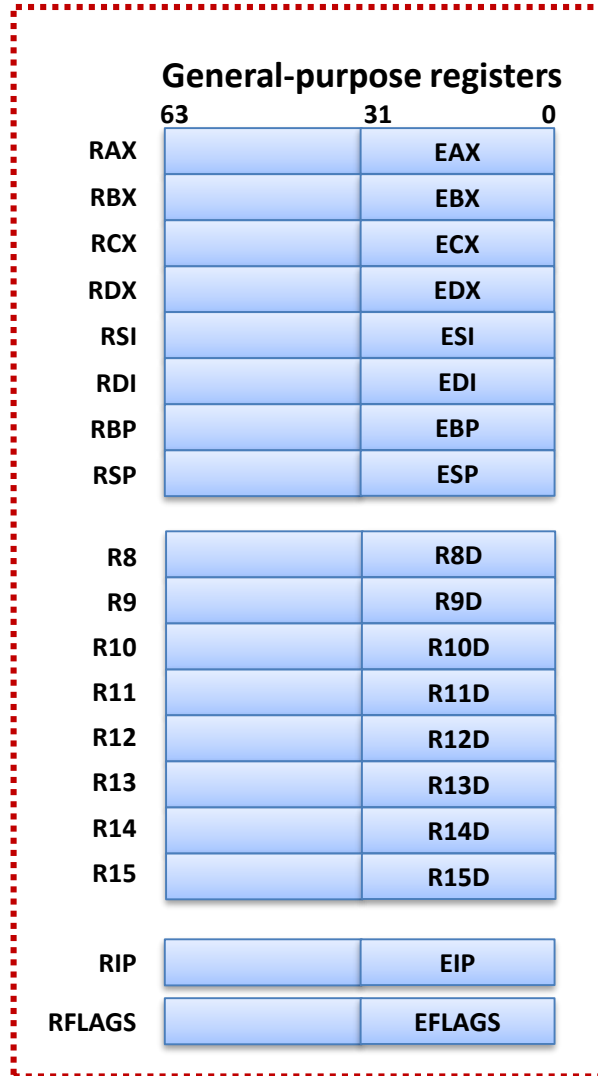
Seoul National University

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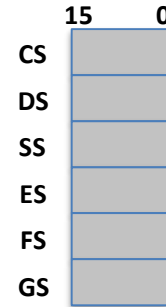
Assembly I: Basic Operations



Basic Execution Environment



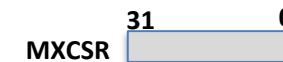
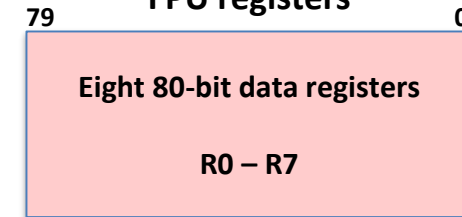
Segment registers



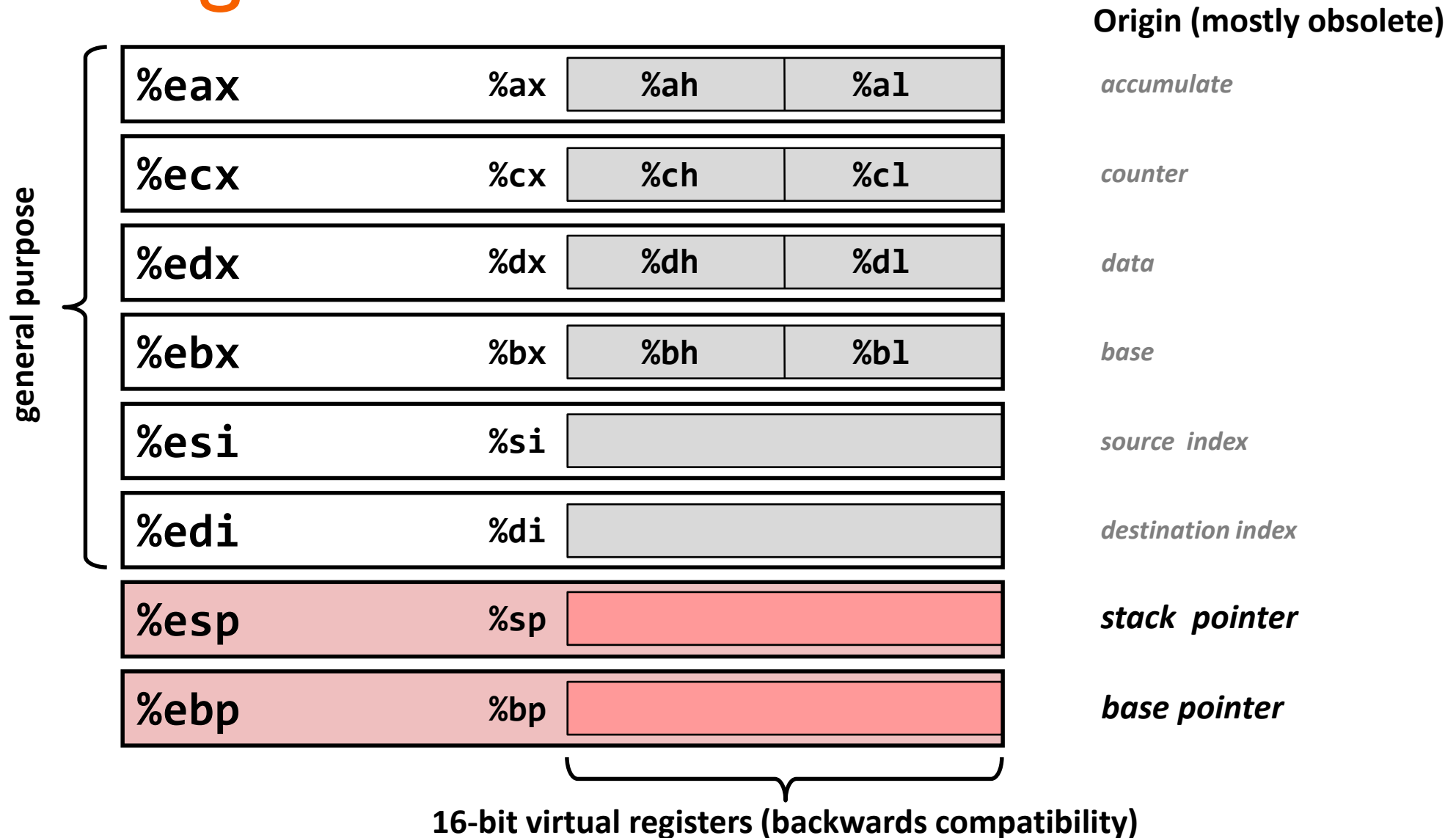
MMX registers



FPU registers



IA-32 Registers



Moving Data (I)

- Moving data: `movq Src, Dest`
 - Move 8-byte (“quad”) word
- Operand types
 - **Immediate**: constant integer data
 - Like C constant, but prefixed with ‘\$’
 - Encoded with 1, 2, or 4 bytes
 - e.g. `$0x400`, `$-533`
 - **Register**: one of 16 integer registers
 - But `%rsp` reserved for special use
 - Others have special uses for particular instructions
 - **Memory**: 8 consecutive bytes of memory
 - Various “addressing modes”

<code>%rax</code>
<code>%rbx</code>
<code>%rcx</code>
<code>%rdx</code>
<code>%rsi</code>
<code>%rdi</code>
<code>%rsp</code>
<code>%rbp</code>
<code>%r8</code>
<code>%r9</code>
<code>%r10</code>
<code>%r11</code>
<code>%r12</code>
<code>%r13</code>
<code>%r14</code>
<code>%r15</code>

Moving Data (2)

- **movq** operand combinations
 - Cannot do memory-memory transfer with a single instruction

	Source	Destination	C Analog	
movq	Imm	Reg	movq \$0x4, %rax	temp = 0x4;
		Mem	movq \$-147, (%rax)	*p = -147;
	Reg	Reg	movq %rax, %rdx	temp2 = temp1;
		Mem	movq %rax, (%rdx)	*p = temp;
	Mem	Reg	movq (%rax), %rdx	temp = *p;

Simple Addressing Modes

- **Normal** **(R)** **Mem[Reg[R]]**
 - Register R specifies memory address
 - Pointer dereferencing in C
 - e.g. `movq (%rcx), %rax`

- **Displacement** **D(R)** **Mem[Reg[R]+D]**
 - Register R specifies start of memory region
 - Constant displacement D specifies offset
 - e.g. `movq 8(%rbp), %rdx`

General Addressing Modes

- **D(Rb, Ri, S)** $\text{Mem}[\text{Reg}[\text{Rb}] + \text{S} * \text{Reg}[\text{Ri}] + \text{D}]$
 - **D:** constant “displacement”: 1, 2, or 4 bytes
 - **Rb:** Base register: any of 16 integer registers
 - **Ri:** Index register: any, except for %rsp
 - **S:** Scale: 1, 2, 4, or 8

- **Special cases**
 - **(Rb, Ri)** $\text{Mem}[\text{Reg}[\text{Rb}] + \text{Reg}[\text{Ri}]]$
 - **D(Rb, Ri)** $\text{Mem}[\text{Reg}[\text{Rb}] + \text{Reg}[\text{Ri}] + \text{D}]$
 - **(Rb, Ri, S)** $\text{Mem}[\text{Reg}[\text{Rb}] + \text{S} * \text{Reg}[\text{Ri}]]$
 - Useful to access arrays and structures

Addressing Modes: Example

- Address computation

%rdx

0xf000

%rcx

0x0100

Expression	Computation	Address
$0x8(\%rdx)$	$0xf000 + 0x8$	0xf008
$(\%rdx, \%rcx)$	$0xf000 + 0x100$	0xf100
$(\%rdx, \%rcx, 4)$	$0xf000 + 4 * 0x100$	0xf400
$0x80(, \%rdx, 2)$	$2 * 0xf000 + 0x80$	0x1e080

Swap Example

- Source code in C:

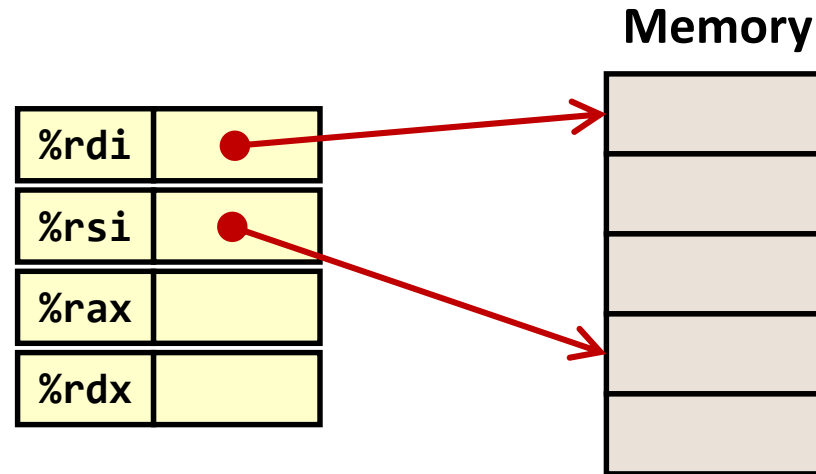
```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

- Corresponding assembly code:

```
swap:
    movq    (%rdi), %rax
    movq    (%rsi), %rdx
    movq    %rdx, (%rdi)
    movq    %rax, (%rsi)
    ret
```

Understanding Swap (I)

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```



Register Allocation (By compiler)

Register	Variable
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)   # *xp = t1
movq    %rax, (%rsi)   # *yp = t0
ret
```

Understanding Swap (2)

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

%rdi	0x120
%rsi	0x100
%rax	
%rdx	

Memory

0x120	123
0x118	
0x110	
0x108	
0x100	456

Register Allocation (By compiler)

Register	Variable
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret
```

Understanding Swap (3)

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

%rdi	0x120
%rsi	0x100
%rax	123
%rdx	

Memory

0x120	123
0x118	
0x110	
0x108	
0x100	456

Register Allocation (By compiler)

Register	Variable
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)   # *xp = t1
movq    %rax, (%rsi)   # *yp = t0
ret
```

Understanding Swap (4)

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

%rdi	0x120
%rsi	0x100
%rax	123
%rdx	456

Memory

0x120	123
0x118	
0x110	
0x108	
0x100	456

Register Allocation (By compiler)

Register	Variable
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

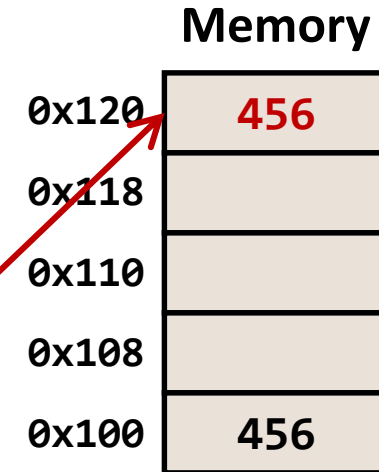
swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret
```

Understanding Swap (5)

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

%rdi	0x120
%rsi	0x100
%rax	123
%rdx	456



Register Allocation (By compiler)

Register	Variable
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)   # *xp = t1
movq    %rax, (%rsi)   # *yp = t0
ret
```

Understanding Swap (5)

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

%rdi	0x120
%rsi	0x100
%rax	123
%rdx	456

Memory

0x120	456
0x118	
0x110	
0x108	
0x100	123

Register Allocation (By compiler)

Register	Variable
%rdi	xp
%rsi	yp
%rax	t0
%rdx	t1

swap:

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)    # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret
```

Arithmetic/Logical Operations (I)

■ Two operand instructions

- **addq** Src, Dest $\text{Dest} = \text{Dest} + \text{Src}$
- **subq** Src, Dest $\text{Dest} = \text{Dest} - \text{Src}$
- **mulq** Src, Dest $\text{Dest} = \text{Dest} * \text{Src}$ (unsigned)
- **imulq** Src, Dest $\text{Dest} = \text{Dest} * \text{Src}$ (signed)
- **salq** Src, Dest $\text{Dest} = \text{Dest} \ll \text{Src}$ (= **shlq**)
- **sarq** Src, Dest $\text{Dest} = \text{Dest} \gg \text{Src}$ (arithmetic)
- **shrq** Src, Dest $\text{Dest} = \text{Dest} \gg \text{Src}$ (logical)
- **xorq** Src, Dest $\text{Dest} = \text{Dest} \wedge \text{Src}$
- **andq** Src, Dest $\text{Dest} = \text{Dest} \& \text{Src}$
- **orq** Src, Dest $\text{Dest} = \text{Dest} | \text{Src}$

Arithmetic/Logical Operations (2)

- One operand instructions

- `incq Dest` $\text{Dest} = \text{Dest} + 1$
- `decq Dest` $\text{Dest} = \text{Dest} - 1$
- `negq Dest` $\text{Dest} = -\text{Dest}$
- `notq Dest` $\text{Dest} = \sim \text{Dest}$

- See books for more instructions

Address Computation Instruction

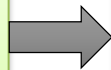
▪ `leaq Src, Dest`

- `Src` is address mode expression
- Set `Dest` to address denoted by expression

▪ Uses

- Computing address without a memory reference
 - e.g. translation of `p = &x[i];`
- Computing arithmetic expression of the form $x + k*y$
 - $k = 1, 2, 4, \text{ or } 8$

```
long m12 (long x) {  
    return x * 12;  
}
```



```
leaq (%rdi, %rdi, 2), %rax  
salq $2, %rax
```

Example: arith

```
long arith (long x, long y, long z) {  
    long t1 = x+y;  
    long t2 = z+t1;  
    long t3 = x+4;  
    long t4 = y * 48;  
    long t5 = t3 + t4;  
    long rval = t2 * t5;  
    return rval;  
}
```

```
x in %rdi  
y in %rsi  
z in %rdx
```

```
arith:  
    leaq (%rdi, %rsi), %rax      # %rax = x + y (t1)  
    addq %rdx, %rax             # %rax = z + t1 (t2)  
    leaq (%rsi, %rsi, 2), %rdx  # %rdx = 3 * y  
    salq $4, %rdx              # %rdx = 48 * y (t4)  
    leaq 4(%rdi, %rdx), %rcx    # %rcx = x + t4 + 4 (t5)  
    imulq %rcx, %rax           # %rax = t5 * t2 (rval)  
    ret
```

Example: logical

```
int logical (int x, int y) {  
    int t1 = x^y;  
    int t2 = t1 >> 17;  
    int mask = (1 << 13) - 7;  
    int rval = t2 & mask;  
    return rval;  
}
```

x in %edi
y in %esi

```
logical:  
    movl %edi, %eax           # %eax = x  
    xorl %esi, %eax          # %eax = x ^ y (t1)  
    sarl $17, %eax           # %eax = t1 >> 17 (t2)  
    andl $8185, %eax         # %eax = t2 & 8185 (rval)  
    ret
```

CISC Properties

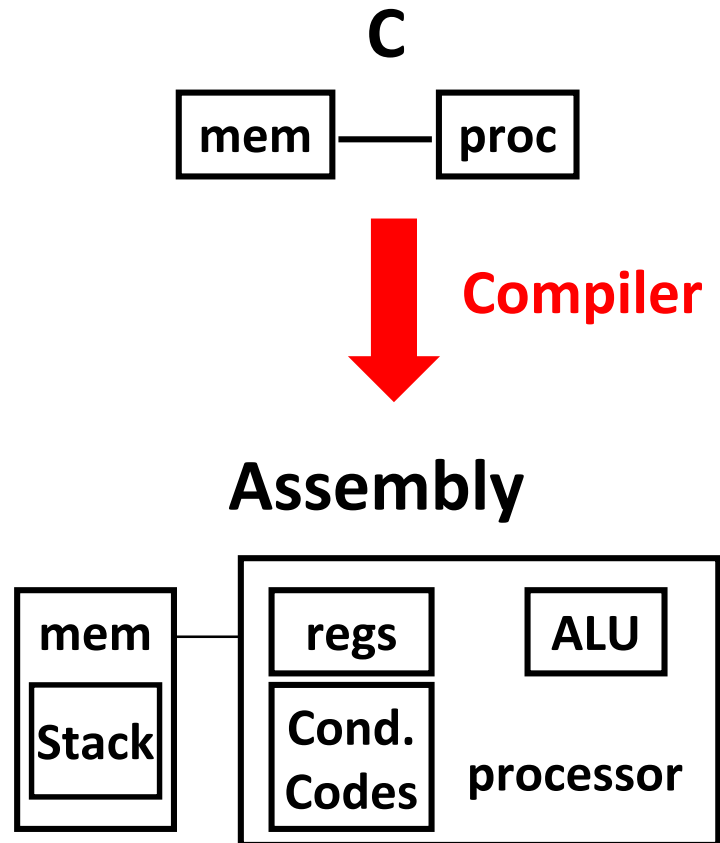
- **Complex Instruction Set Computer (CISC)**
 - Instruction can reference different operand types
 - Immediate, register, memory
 - Arithmetic operations can read/write memory
 - Source or destination can be a memory address
 - Memory reference can involve complex computation
 - $R_b + S * R_i + D$
 - Useful for arithmetic expressions, too
 - Instructions can have varying lengths
 - x86-64 instructions can range from 1 to 15 bytes

Machine-level Programming

- Assembly code is textual form of binary object code
- Low-level representation of program
 - Explicit manipulation of registers
 - Simple and explicit instructions
 - Minimal concept of data types
 - Many C control constructs must be implemented with multiple instructions

Summary

Machine Models



Data

- 1) char
- 2) int, float
- 3) double
- 4) struct, array
- 5) pointer

Control

- 1) loops
- 2) conditionals
- 3) switch
- 4) Proc. call
- 5) Proc. return

- 1) byte
- 2) 2-byte word
- 3) 4-byte long word
- 4) 8-byte quad word
- 5) contiguous byte allocation
- 6) address of initial byte

- 1) branch/jump
- 2) call
- 3) ret