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RISC-V User-level ISA Cheat Sheet



The RISC-V Instruction Set

- A completely open ISA that is freely available to academia and industry
- Fifth RISC ISA design developed at UC Berkeley
 - RISC-I (1981), RISC-II (1983), SOAR (1984), SPUR (1989), and RISC-V (2010)
- Now managed by the RISC-V Foundation (<http://riscv.org>)
- Typical of many modern ISAs
 - See RISC-V Reference Card (or Green Card)
- Similar ISAs have a large share of embedded core market
 - Applications in consumer electronics, network/storage equipment, cameras, printers, ...

RISC-V ISAs

- Three base integer ISAs, one per address width
 - RV32I, RV64I, RV128I
 - RV32I: Only 40 instructions defined
 - RV32E: Reduced version of RV32I with 16 registers for embedded systems
- Standard extensions
 - Standard RISC encoding in a fixed 32-bit instruction format
 - C extension offers shorter 16-bit versions of common 32-bit RISC-V instructions (can be intermixed with 32-bit instructions)

Name	Extension
M	Integer Multiply/Divide
A	Atomic Instructions
F	Single-precision FP
D	Double-precision FP
G	General-purpose (= IMAFD)
Q	Quad-precision FP
C	Compressed Instructions

RISC-V Registers

#	Name	Usage
x0	zero	Hard-wired zero
x1	ra	Return address
x2	sp	Stack pointer
x3	gp	Global pointer
x4	tp	Thread pointer
x5	t0	Temporaries (Caller-save registers)
x6	t1	
x7	t2	
x8	s0/fp	Saved register / Frame pointer
x9	s1	Saved register
x10	a0	Function arguments / Return values
x11	a1	
x12	a2	Function arguments
x13	a3	
x14	a4	
x15	a5	

#	Name	Usage
x16	a6	Function arguments
x17	a7	
x18	s2	Saved registers (Callee-save registers)
x19	s3	
x20	s4	
x21	s5	
x22	s6	
x23	s7	
x24	s8	
x25	s9	
x26	s10	
x27	s11	
x28	t3	Temporaries (Caller-save registers)
x29	t4	
x30	t5	
x31	t6	
	pc	Program counter

Arithmetic Operations

Instruction	Type	Example	Meaning
Add	R	add rd, rs1, rs2	$R[rd] = R[rs1] + R[rs2]$
Subtract	R	sub rd, rs1, rs2	$R[rd] = R[rs1] - R[rs2]$
Add immediate	I	addi rd, rs1, imm12	$R[rd] = R[rs1] + \text{SignExt}(imm12)$
Set less than	R	slt rd, rs1, rs2	$R[rd] = (R[rs1] < R[rs2])? 1 : 0$
Set less than immediate	I	slti rd, rs1, imm12	$R[rd] = (R[rs1] < \text{SignExt}(imm12))? 1 : 0$
Set less than unsigned	R	sltu rd, rs1, rs2	$R[rd] = (R[rs1] <_u R[rs2])? 1 : 0$
Set less than immediate unsigned	I	sltiu rd, rs1, imm12	$R[rd] = (R[rs1] <_u \text{SignExt}(imm12))? 1 : 0$
Load upper immediate	U	lui rd, imm20	$R[rd] = \text{SignExt}(imm20 \ll 12)$
Add upper immediate to PC	U	auipc rd, imm20	$R[rd] = PC + \text{SignExt}(imm20 \ll 12)$

Logical Operations

Instruction	Type	Example	Meaning
AND	R	and rd, rs1, rs2	$R[rd] = R[rs1] \& R[rs2]$
OR	R	or rd, rs1, rs2	$R[rd] = R[rs1] R[rs2]$
XOR	R	xor rd, rs1, rs2	$R[rd] = R[rs1] \wedge R[rs2]$
AND immediate	I	andi rd, rs1, imm12	$R[rd] = R[rs1] \& \text{SignExt}(imm12)$
OR immediate	I	ori rd, rs1, imm12	$R[rd] = R[rs1] \text{SignExt}(imm12)$
XOR immediate	I	xori rd, rs1, imm12	$R[rd] = R[rs1] \wedge \text{SignExt}(imm12)$
Shift left logical	R	sll rd, rs1, rs2	$R[rd] = R[rs1] \ll R[rs2]$
Shift right logical	R	srl rd, rs1, rs2	$R[rd] = R[rs1] \gg R[rs2]$ (logical)
Shift right arithmetic	R	sra rd, rs1, rs2	$R[rd] = R[rs1] \gg R[rs2]$ (arithmetic)
Shift left logical immediate	I	slli rd, rs1, shamt	$R[rd] = R[rs1] \ll shamt$
Shift right logical imm.	I	srlr rd, rs1, shamt	$R[rd] = R[rs1] \gg shamt$ (logical)
Shift right arithmetic immediate	I	srair rd, rs1, shamt	$R[rd] = R[rs1] \gg shamt$ (arithmetic)

Data Transfer Operations

Instruction	Type	Example	Meaning
Load doubleword	I	ld rd, imm12(rs1)	$R[rd] = \text{Mem}_8[R[rs1] + \text{SignExt}(\text{imm12})]$
Load word	I	lw rd, imm12(rs1)	$R[rd] = \text{SignExt}(\text{Mem}_4[R[rs1] + \text{SignExt}(\text{imm12})])$
Load halfword	I	lh rd, imm12(rs1)	$R[rd] = \text{SignExt}(\text{Mem}_2[R[rs1] + \text{SignExt}(\text{imm12})])$
Load byte	I	lb rd, imm12(rs1)	$R[rd] = \text{SignExt}(\text{Mem}_1[R[rs1] + \text{SignExt}(\text{imm12})])$
Load word unsigned	I	lwu rd, imm12(rs1)	$R[rd] = \text{ZeroExt}(\text{Mem}_4[R[rs1] + \text{SignExt}(\text{imm12})])$
Load halfword unsigned	I	lhu rd, imm12(rs1)	$R[rd] = \text{ZeroExt}(\text{Mem}_2[R[rs1] + \text{SignExt}(\text{imm12})])$
Load byte unsigned	I	lbu rd, imm12(rs1)	$R[rd] = \text{ZeroExt}(\text{Mem}_1[R[rs1] + \text{SignExt}(\text{imm12})])$
Store doubleword	S	sd rs2, imm12(rs1)	$\text{Mem}_8[R[rs1] + \text{SignExt}(\text{imm12})] = R[rs2]$
Store word	S	sw rs2, imm12(rs1)	$\text{Mem}_4[R[rs1] + \text{SignExt}(\text{imm12})] = R[rs2](31:0)$
Store halfword	S	sh rs2, imm12(rs1)	$\text{Mem}_2[R[rs1] + \text{SignExt}(\text{imm12})] = R[rs2](15:0)$
Store byte	S	sb rs2, imm12(rs1)	$\text{Mem}_1[R[rs1] + \text{SignExt}(\text{imm12})] = R[rs2](7:0)$

Control Transfer Instructions

Instruction	Type	Example	Meaning
Branch equal	SB	beq rs1, rs2, imm12	if (R[rs1] == R[rs2]) pc = pc + SignExt(imm12 << 1)
Branch not equal	SB	bne rs1, rs2, imm12	if (R[rs1] != R[rs2]) pc = pc + SignExt(imm12 << 1)
Branch greater than or equal	SB	bge rs1, rs2, imm12	if (R[rs1] >= R[rs2]) pc = pc + SignExt(imm12 << 1)
Branch greater than or equal unsigned	SB	bgeu rs1, rs2, imm12	if (R[rs1] >= _u R[rs2]) pc = pc + SignExt(imm12 << 1)
Branch less than	SB	blt rs1, rs2, imm12	if (R[rs1] < R[rs2]) pc = pc + SignExt(imm12 << 1)
Branch less than unsigned	SB	bltu rs1, rs2, imm12	if (R[rs1] < _u R[rs2]) pc = pc + SignExt(imm12 << 1)
Jump and link	UJ	jal rd, imm20	R[rd] = PC + 4 PC = PC + SignExt(imm20 << 1)
Jump and link register	I	jalr rd, imm12(rs1)	R[rd] = PC + 4 PC = (R[rs1] + SignExt(imm12)) & (~1)

Assembler Pseudo-Instructions

Pseudo-instruction	Base instruction(s)	Meaning
<code>li rd, imm</code>	<code>addi rd, x0, imm</code>	Load immediate
<code>la rd, symbol</code>	<code>auipc rd, D[31:12]+D[11]</code> <code>addi rd, rd, D[11:0]</code>	Load absolute address where $D = \text{symbol} - \text{pc}$
<code>mv rd, rs</code>	<code>addi rd, rs, 0</code>	Copy register
<code>not rd, rs</code>	<code>xori rd, rs, -1</code>	One's complement
<code>neg rd, rs</code>	<code>sub rd, x0, rs</code>	Two's complement
<code>bgt{u} rs, rt, offset</code>	<code>blt{u} rt, rs, offset</code>	Branch if $>$ (u: unsigned)
<code>ble{u} rs, rt, offset</code>	<code>bge{u} rt, rs, offset</code>	Branch if \geq (u: unsigned)
<code>b{eq ne}z rs, offset</code>	<code>b{eq ne} rs, x0, offset</code>	Branch if $\{ = \neq \}$
<code>b{ge lt}z rs, offset</code>	<code>b{ge lt} rs, x0, offset</code>	Branch if $\{ \geq < \}$
<code>b{le gt}z rs, offset</code>	<code>b{ge lt} x0, rs, offset</code>	Branch if $\{ \leq > \}$
<code>j offset</code>	<code>jal x0, offset</code>	Unconditional jump
<code>call offset</code>	<code>jal ra, offset</code>	Call subroutine (near)
<code>ret</code>	<code>jalr x0, 0(ra)</code>	Return from subroutine
<code>nop</code>	<code>addi x0, x0, 0</code>	No operation