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Operating Systems
Lab. I



What is system call?

- User space applications run with restricted privilege
- They have to request some operations to the OS
- System call is an exception that intentionally made by application for this purpose

3 modes of RISC-V

Machine mode

- Have full privilege
- CPU starts in machine mode

Supervisor mode

- Allowed to execute privileged instructions
 - Enabling and disabling interrupts
 - Reading and writing the register that holds address of page table
 - **–** ...
- The kernel runs in supervisor mode

User mode

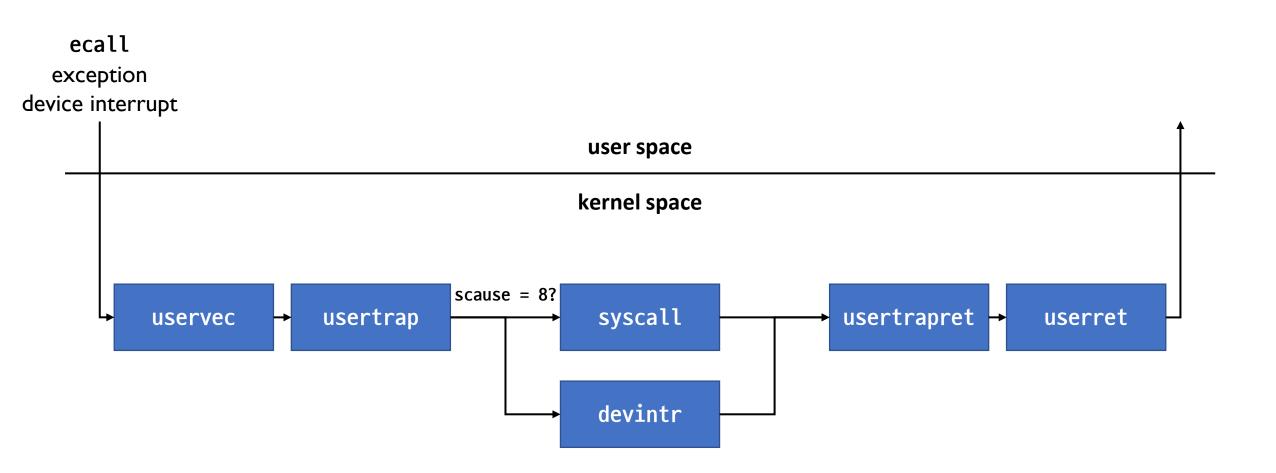
• User space applications runs in user mode

ecall instruction

- In RISC-V an application can use "ecall" instruction to invoke system call
- e.g.) Fork function in userspace

```
fork:
li a7, SYS_fork
ecall
ret
```

Traps from user space



Special RISC-V registers

- stvec: The address of trap handler
- sepc: Register to save program counter when a trap occurs
- scause: The reason of a trap
- sscratch: The address of trapframe

satp: Current page table

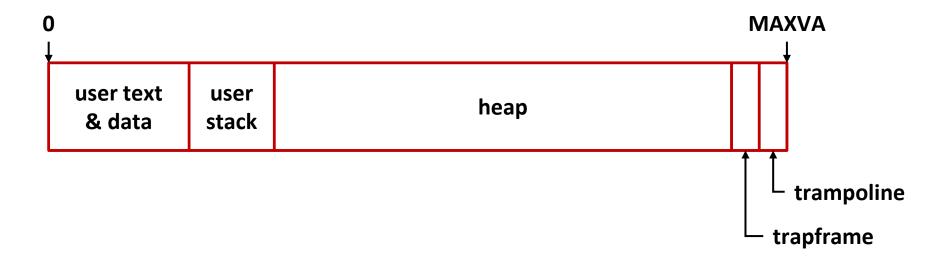
On trap...

- CPU runs trap handler saved in Stvec register
- Program counter is saved to Sepc register
- The reason of the trap is saved to scause register

uservec

- Saves all register values to trapframe
- Change satp(page table register) to kernel page table
- uservec is located on trampoline page because changing page table is needed

Virtual Address Space Layout



usertrap

- usertrap determines the cause of the trap, and handles it
- Set stvec to kernelvec
- Call syscall if the trap was caused by system call, devintr if the trap was caused by device interrupt

syscall

- Gets system call number from a7 register saved in trapframe
- Calls system call handler
 - e.g.) sys_fork function in kernel space
- Saves return value to a0 register in trapframe

usertrapret

- Saves kernel page table, kernel stack to trapframe (To be used for next trap from user space)
- Restore stvec to refer uservec
- Restore sepc to previously saved user program counter

userret

- Switches satp to process's user page table
- Restores all register values from trapframe
- userret is also located on trampoline page because changing page table is needed

Project #2 – System call

- Your job is to implement process group and two system calls
- Why process group is needed?
 - e.g.) tar -c file1 file2 | xz --threads=4 > files.tar.xz
 - If you want to abort the job, two processes(tar and xz) must be interrupted together
- You have to implement:
 - setpgid to set process group ID
 - getpgid to get process group ID

setpgid system call

- int setpgid(int pid, int pgid);
- setpgid should change the process group ID of process specified by pid to pgid
- pid and pgid are non-negative integers
- If pid is 0, it should change process group ID of the calling process(process that invoked the system call)
- If pgid is 0, it should change process group ID same as pid

- On success, return 0
- On failure, return I

getpgid system call

- int getpgid(int pid);
- getpgid should return process group ID of process specified by pid
- If pid is 0, it should return process group ID of the calling process
- On success, return pgid
- On failure, return I

Displaying process list

- xv6 prints list of processes when Ctrl+P is pressed
- You have to modify this to print pgid
 - format: pid pgid state process_name

```
xv6 kernel is booting
hart 1 starting
hart 2 starting
init: starting sh
$
1 1 sleep init
2 1 sleep sh
```

You may want to see...

- defs.h
 - For function definitions
- proc.h, proc.c
 - For process related functions
- console.c
 - For console input handling
- syscall.c, sysproc.c
 - For system call implementation

When you do your project,

- Please only modify Makefile and files in kernel directory
 - Please fill your student id to STUDENTID variable in Makefile
 - Modifications to user directory will be ignored by grading script
- Please remove all the debugging outputs before you submit

- Please read the project description carefully
 - https://github.com/snu-csl/os-pa2
- Skeleton code is on xv6-riscv-snu repository's pa2 branch
 - git clone https://github.com/snu-csl/xv6-riscv-snu
 - git checkout pa2
- Archive your source to tarball and submit it to sys server
 - Type "make submit" to archive your source
 - And upload it to https://sys.snu.ac.kr

Thank you!

If you have any questions, feel free to ask us in KakaoTalk